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## Key Mismatch Attack on ThreeBears, Frodo and Round5

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## Outline

1. Targeted schemes
2. Key Mismatch Attack
3. State-of-the-art
4. Our attack
5. Results
6. Conclusion

## Targeted schemes

- **ThreeBears**

- based on Integer Module Learning with Errors (I-MLWE)
- NIST round 2 candidate
- *C. Gu: “Integer Version of Ring-LWE and its Applications”, 2017*

- **Frodo**

- based on Learning with Errors (LWE)
- NIST alternative round 3 candidate
- *J. Bos et al.: “Frodo: Take off the ring! Practical, Quantum-Secure Key Exchange from LWE”, 2016*

- **Round5**

- based on Learning with Rounding (LWR) and Ring Learning with Rounding (RLWR)
- NIST round 2 candidate
- *A. Banerjee et al.: “Pseudorandom Functions and Lattices”, 2011*

## Key Mismatch Oracle Attack

- Consider some Public Key Encryption (PKE) with a fixed secret key  $sk$
- The goal of the attacker is to recover  $sk$  using the key mismatch oracle:

- INPUT: - arbitrarily chosen ciphertext  $ct$  (not necessarily computed according to the specification)  
- arbitrary plaintext  $pt$
- OUTPUT: 
$$\begin{cases} + & \text{if } \text{decryption}(sk, ct) = pt \\ - & \text{if } \text{decryption}(sk, ct) \neq pt \end{cases}$$

## Practical relevance

- Access to the Key Mismatch Oracle even for actively secure (CCA) variants using e.g. side-channel attacks
- There is a risk of key reuse even though it is forbidden by the specification
- Significant state-of-the-art on the topic, e.g.:
  - S. Fluhrer: “*Cryptanalysis of ring-LWE based key exchange with key share reuse*”, 2016
  - S. Vaudenay et al.: “*Misuse Attacks on Post-Quantum Cryptosystems*”, EUROCRYPT 2019
  - S. Vaudenay et al.: “*Classical Misuse Attacks on NIST Round 2 PQC: The Power of Rank-Based Schemes*”, ACNS 2020
  - P. Ravi et al.: “*Generic Side-channel attacks on CCA-secure lattice-based PKE and KEM schemes*”, CHES 2020
  - S. Okada et al.: “*Improving Key Mismatch Attack on NewHope with Fewer Queries*”, ACISP 2020

## Key mismatch oracle attack in the prior art

- Previous attacks target secret coefficients one by one
- Common technique – consider only queries such that:
  - the possible mismatch between the decrypted plaintext and the chosen plaintext can happen only on one position
  - the bit on this position depends only on the targeted secret coefficient
- Differences – how the output from the oracle is utilized:
  - “favorable cases” (ACISP 2020)
  - recover linear equations with the secret key as unknown (EUROCRYPT 2019)
  - oracle output tells if a coefficient is greater than a given threshold (ACNS 2020)
  - associate output sequences with targeted coefficients (CHES 2020)

## Idea of our attack

- Secret coefficients targeted in tuples, not necessary one by one
- Gradually reduce the possibilities for the targeted tuple

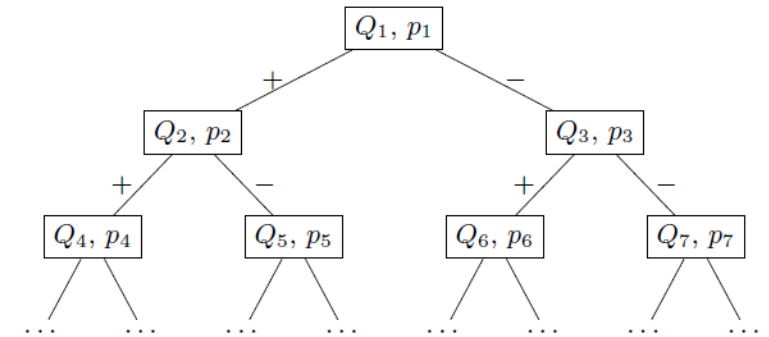
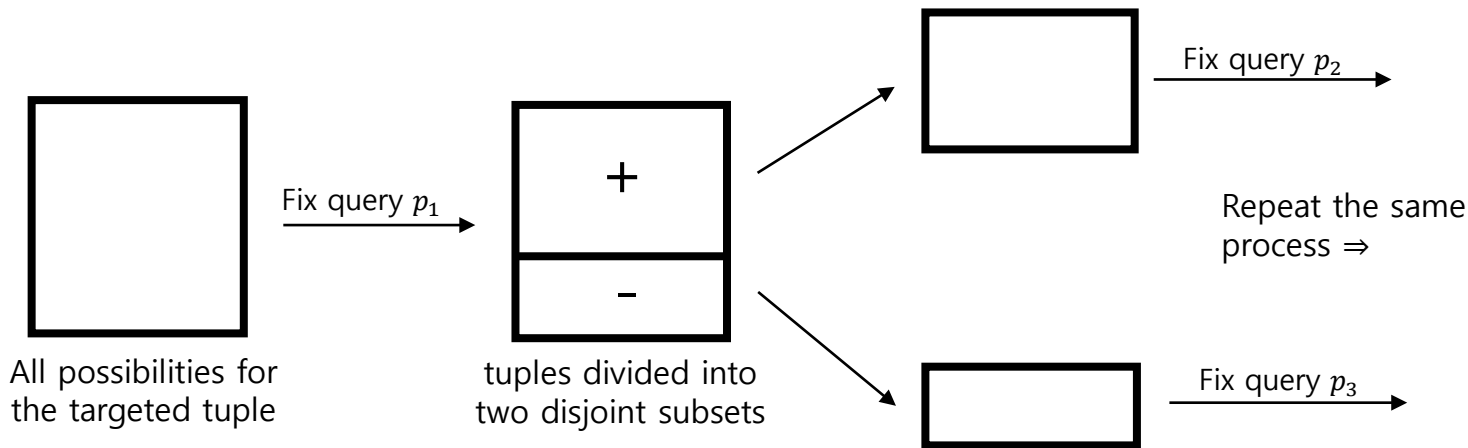


Fig. 2. Tree structure.

❖ We want  $|Q_i| = 1$  for the leaves

## The attack

- The attacker follows a path from the root to some leaf according to the outputs from the oracle
- The attacker does not perform any computation, all the queries are stored within the tree
- The number of queries to recover some tuple equals the depth of the leaf corresponding to this tuples

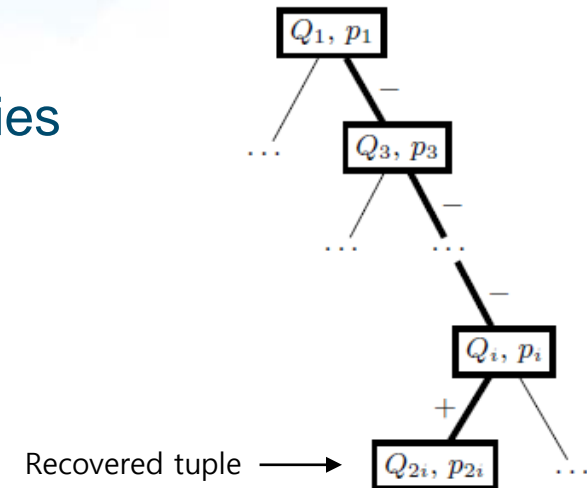


Fig. 3. Path in the tree.



## Construction of trees

- Tree is constructed recursively: looking for the tree which minimizes the expected number of queries to the oracle
- The expected number of queries: weighted average of the probabilities of the tuples and of the depths of the leaves corresponding to these tuples
- Not possible to try each tree  $\Rightarrow$  the following heuristic is used: split a set of possible tuples such that the two disjoint subsets have similar probabilities

# Results for ThreeBears

- We provide the first attack on ThreeBears
- Coefficients targeted only one by one

Error-correcting code	NIST security level	Expected number of queries	Success probability
Yes	1	1 414	100%
Yes	3	1 638	100%
Yes	5	2 223	100%
No	1	1 443	100%
No	3	2 150	100%
No	5	2 847	100%

# Results for Frodo

- Coefficients targeted one by one and by pairs (called dimension of the attack)
- Existing attack by Vaudenay et al. from EUROCRYPT 2019

	NIST security level	Dimension of the attack	Expected number of queries	Success probability
EUROCRYPT 2019	1	-	65 536	not clear
	1	1	18 359	100%
	1	2	18 239	100%
	3	1	25 934	100%
	3	2	25 672	100%
	5	1	29 377	100%
	5	2	28 008	100%

# Results for Round5

- Coefficients targeted one by one, by pairs, triplets and quadruplets
- Existing attack by Ravi et al. from CHES 2020

	variant	NIST security level	Dimension of the attack	Expected number of queries	Success probability
CHES 2020	RLWR+ECC	1	-	978	100%
	RLWR+ECC	1	4	656	100%
	RLWR+ECC	3	4	1277	100%
	RLWR	1	3	687	100%
	RLWR	3	2	1221	100%
	LWR	1	4	5 790	100%
	LWR	3	4	8 436	100%

## Conclusion

- The first key mismatch attack on ThreeBears and variants of Round5
- Improved key mismatch attack on Frodo and variant of Round5
- The method is applicable against other LWE-based candidates, e.g. against Kyber, Saber, NewHope
- Targeting bigger tuples (if possible) gives better results, but it is not possible to target arbitrary tuples

**Thank you for your attention!**

**Questions?**